

Measurement-device-independent Randomness from local entangled states

ISCQI-2016
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Figure: Anubhav Chaturvedi



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Link: [Euro. Phys. Lett. 112, 30003 \(2015\)](#)
(doi: [10.1209/0295-5075/112/30003](https://doi.org/10.1209/0295-5075/112/30003))

★ Random numbers have many practical uses in modern science.

- Cryptography
- Statistical research
- Numerical Simulation (eg. Monte Carlo method)
- Lotteries and gambling
- PIN number generation
- Mobile prepaid systems

★ Consider a sequence of two numbers '0' and '1':

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with $p(0) = p(1) = \frac{1}{2}$

Randomness Certification...

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Certification: Mathematically impossible...

★ Using algorithmic information theory it can be shown that true randomness can not exist from a mathematical point of view.

{Chaitin G. J., IBM J. Res. Dev., 21 (1977) 350; Knuth D., The Art of Computer Programming, Semi-numerical Algorithms}



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★ Any one who considers arithmetical methods of producing random digits is, of course, in a state of sin.

——John von Neumann

★ Therefore generation of randomness must rely on unpredictability of physical phenomena, like

Certification: Impossible in classical physics...

★ Coin tossing...



Certification: Impossible in classical physics...

★ Dice rolling...



Certification: Impossible in classical physics...

- ★ However all classical processes ('**coin tossing**'/'**dice rolling**') are deterministic from fundamental point of view.
- ★ This is because, fate of any classical object at any time is completely predictable in Newtonian dynamics.
- ★ Therefore no classical process can be a source of "true" randomness.



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- ★ Therefore no classical process can be a source of "true" randomness.



- ★ So we shift our attention from classical world (CW) to quantum world (QW).

Certification: Possible in QW (?)...

★ Quantum theory:

- **State of a system: Vectors, $|\psi\rangle \in \mathcal{H}$**
- **Observables: Hermitian operator, $\mathcal{A} \in \mathcal{B}(\mathcal{H})$, acting on \mathcal{H}**
$$\mathcal{A} = \sum_i a_i |\alpha_i\rangle\langle\alpha_i|$$
- **Possible measurement results: Eigenvalues of the given observable**
- **Outcome probability: $p(a_1) = |\langle\psi|\alpha_1\rangle|^2$ (Born rule)**

★ Due to the Born's rule, in QW we can obtain our desired randomness.

Certification: Possible in QW (?) {an example}...

★ Consider a spin-1/2 system ($\mathcal{H} \equiv \mathbb{C}^2$)

- **State:** 'up' eigenstate ($|\uparrow\rangle$) of the Pauli σ_z
- **Measurement:** Pauli σ_x observable, whose eigenstate are denoted as $|\leftrightarrow\rangle = \frac{1}{\sqrt{2}}(|\uparrow\rangle \pm |\downarrow\rangle)$.
- **Outcome probability:** $p(|\rightarrow\rangle) = \frac{1}{2}$ and also $p(|\leftarrow\rangle) = \frac{1}{2}$

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★ If we associate '0' ('1') with $|\uparrow\rangle$ ($|\downarrow\rangle$), then we obtain a sequence of '0' and '1' with $p(0) = p(1) = \frac{1}{2}$ and there will be no pattern in the sequence \implies **Randomness**



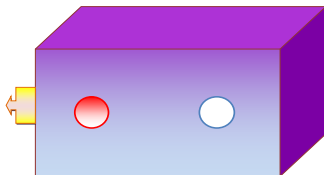
QRNG...

★ Such QRNG already exists:



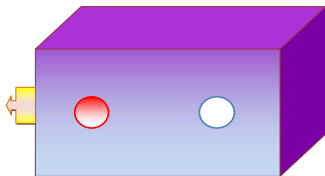
DI Certification...

★ Consider that we are ignored about the internal working of the device, i.e., the device is like a black box with just input and output.



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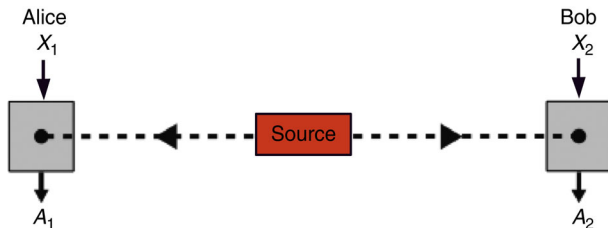
★ In this situation, is it still possible to be certain that the device is producing the outcomes without following any particular pattern?



★ In a recent result it has been shown that DI randomness certification is possible.

★ "Random Numbers Certified by Bell's Theorem", S. Pironio, A. Acin, S. Massar, A. Boyer de la Giroday, D. N. Matsukevich, P. Maunz, S. Olmschenk, D. Hayes, L. Luo, T. A. Manning, C. Monroe, [Nature 464, 1021 \(2010\)](#)

Bell's theorem...



★ **Local Realism (LR):**

$$P(A_1, A_2 | X_1, X_2) = \sum_{\lambda \in \Lambda} \rho(\lambda) P(A_1 | X_1, \lambda) P(A_2 | X_2, \lambda)$$

★ **Bell inequality:** $|\langle X_1 X_2 \rangle + \langle X_1 X'_2 \rangle + \langle X'_1 X_2 \rangle - \langle X'_1 X'_2 \rangle| \leq 2$

★ **Bell inequality can also be derived under two operational assumptions, namely 'predictability' and 'signal-locality'**

Bell's theorem...

★ **Determinism** \wedge **Locality** \Rightarrow **factorizability** \Rightarrow **Bell's inequality (BI)**

i.e., $P(A_1, A_2|X_1, X_2, \psi) = \int_{\lambda \in \Lambda} \mu(\lambda|\psi) P(A_1|X_1, \psi, \lambda) P(A_2|X_2, \psi, \lambda) d\lambda$

- **Determinism (D)** $\Rightarrow P(A_1, A_2|X_1, X_2, \psi, \lambda) \in \{0, 1\}$
- **Locality (L)** $\Rightarrow P(A_1|X_1, X_2, \psi, \lambda) = P(A_1|X_1, \psi, \lambda)$
 $P(A_2|X_1, X_2, \psi, \lambda) = P(A_2|X_2, \psi, \lambda)$
- **Proof:**

$$\begin{aligned} P(A_1, A_2|X_1, X_2, \psi, \lambda) &= P(A_1|A_2, X_1, X_2, \psi, \lambda) P(A_2|X_1, X_2, \psi, \lambda) \\ &= P(A_1|X_1, X_2, \psi, \lambda) P(A_2|X_1, X_2, \psi, \lambda); [D] \\ &= P(A_1|X_1, \psi, \lambda) P(A_2|X_2, \psi, \lambda); [L] \end{aligned}$$

★ **Predictability** \wedge **Signal Locality** \Rightarrow **factorizability** \Rightarrow **BI**

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- **Predictability (P)** $\Rightarrow P(A_1, A_2 | X_1, X_2, \psi) \in \{0, 1\}$
- **Signal Locality (SL)** $\Rightarrow P(A_1 | X_1, X_2, \psi) = P(A_1 | X_1, \psi)$
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Bell's theorem...

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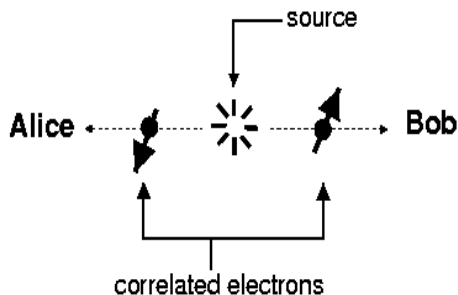
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★ \neg **BI** \wedge **Signal Locality** $\Rightarrow \neg$ **Predictability**

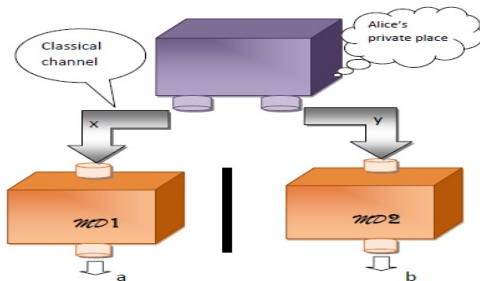
Bell's theorem...

★ Quantum correlation violates BI:



★ Using nonlocal correlation DI randomness certification is possible.

DI Randomness Certification.....



★ Randomness associated with $\{P(ab|xy)\}$, quantified as $H_\infty = -\log_2 \max_{a,b} P(ab|xy)$.

★ Which physical correlation shows this nonlocal properties?

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★ **Entanglement:**

- **Bipartite quantum system** $\rightarrow \mathcal{H}_A \otimes \mathcal{H}_B$
- **Product state:** $|\psi\rangle_A \otimes |\psi\rangle_B \in \mathcal{H}_A \otimes \mathcal{H}_B$
- **Non product States are called entangled:**
$$|\psi\rangle_{AB} = \sum_i c_i |\psi^i\rangle_A \otimes |\psi^i\rangle_B$$
- **Separable states:** $\rho_{AB} = \sum_i p_i \sigma_A^i \otimes \sigma_B^i$; with
 $p_i \geq 0$ & $\sum \liminf_i p_i = 1$

DI Randomness Certification.....

★ Example: Werner class of states

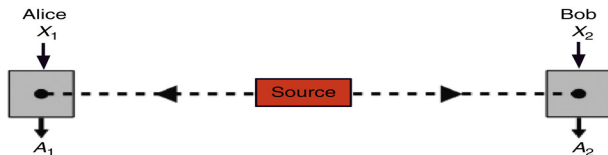
$$W_p = p|\psi^-\rangle\langle\psi^-| + (1-p)\frac{\mathbb{I}}{2} \otimes \frac{\mathbb{I}}{2}$$

with $0 \leq p \leq 1$ and $|\psi^-\rangle = \frac{1}{\sqrt{2}}(|\uparrow\rangle|\downarrow\rangle - |\downarrow\rangle|\uparrow\rangle)$

- $p > \frac{1}{\sqrt{2}}$: Violates BI (useful for DI certification)
- $p \leq \frac{1}{2}$: LHV for PV
- $p \leq \frac{5}{12}$: LHV for POVM
- $p \geq \frac{1}{3}$: Entangled

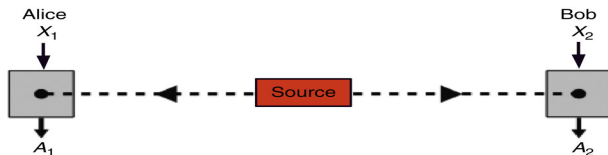
★ Local entangled states are not useful for DI randomness certification

★ Semi-quantum game (F. Buscemi):



- Instead of classical inputs, quantum states $\{|\phi^x\rangle_{\alpha'}\}_{x \in X}$ and $\{|\psi^y\rangle_{\beta'}\}_{y \in Y}$, chosen from Hilbert spaces $\mathcal{H}_{\alpha'}$ and $\mathcal{H}_{\beta'}$, respectively, are sent
- For every entangled state quantum inputs can be chosen in such a way that the produced correlation cannot be achieved by local operation and shared randomness

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MDI Randomness Certification.....

- We consider the following particular semi-quantum game:
- The input quantum states are chosen from a regular tetrahedron on the Bloch sphere i.e.,

$$|\phi^x\rangle\langle\phi^x| = \frac{\mathbb{I} + \vec{v}_x \cdot \vec{\sigma}}{2}, \quad |\psi^y\rangle\langle\psi^y| = \frac{\mathbb{I} + \vec{v}_y \cdot \vec{\sigma}}{2},$$

for $x, y = 1, \dots, 4$ we have $\vec{v}_1 = \frac{(1,1,1)}{\sqrt{3}}$, $\vec{v}_2 = \frac{(1,-1,-1)}{\sqrt{3}}$, $\vec{v}_3 = \frac{(-1,1,-1)}{\sqrt{3}}$ and $\vec{v}_4 = \frac{(1,-1,1)}{\sqrt{3}}$; and $\vec{\sigma} = (\sigma_1, \sigma_2, \sigma_3)$ with σ_i ($i = 1, 2, 3$) being the Pauli matrices

- The POVM $\{\mathcal{M}_a^{\alpha'\alpha}\}_{a \in \{0,1\}}$ is given by

$$\mathcal{M}_1^{\alpha'\alpha} = |\phi^+\rangle\langle\phi^+|, \quad \mathcal{M}_0^{\alpha'\alpha} = \mathbb{I} - |\phi^+\rangle\langle\phi^+|,$$

MDI Randomness Certification.....

- Using the above quantum-input classical-output statistics one can construct the following MDI-entanglement witness (Branciard et al.):

$$I(P) = \frac{5}{8} \sum_{x=y} p(1, 1 | |\phi^x\rangle, |\psi^y\rangle) - \frac{1}{8} \sum_{x \neq y} p(1, 1 | |\phi^x\rangle, |\psi^y\rangle).$$

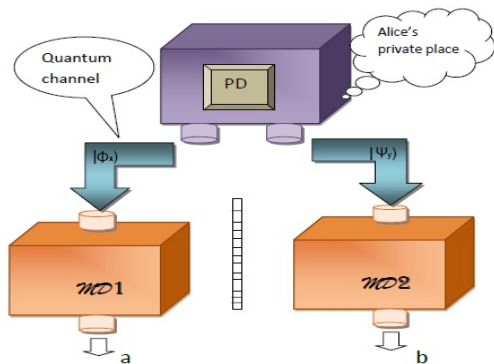
Here P denotes the probability distribution

$$\{p(a, b | |\phi^x\rangle, |\psi^y\rangle) | a, b = 0, 1; x, y = 1, \dots, 4\}.$$

- For the Werner states the above expression becomes $I(P_{\rho^v}) = \frac{1-3v}{16}$, which is negative for $v > \frac{1}{3}$. For any separable state ρ , $I(P_\rho) = 0$

MDI Randomness Certification.....

★ Measurement-DI (MDI) randomness certification:



★ MDI min-entropic source

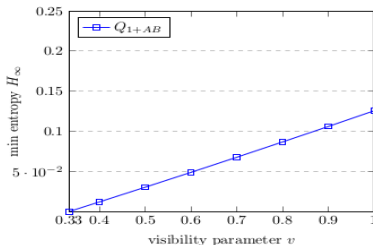
To find the minimum randomness associated with the probability distribution $P = \{p(ab|xy)\}$ one has to solve the following optimization problem,

$$\begin{aligned} p^*(ab|xy) &= \max p(ab|xy) \\ \text{subject to } I(P) &= \frac{1 - 3v}{16} \\ p(ab|xy) &\in Q, \end{aligned}$$

- where $I(P) = \frac{5}{8} \sum_{x=y} p(1, 1 | |\phi^x\rangle, |\psi^y\rangle) - \frac{1}{8} \sum_{x \neq y} p(1, 1 | |\phi^x\rangle, |\psi^y\rangle)$
- the minimum random bits is therefore $H_\infty(AB|XY) = -\log_2 \max_{ab} p_q^*(ab|xy)$.
- While the optimization problem is computationally tough, we solve for a relaxed condition $p(ab|xy) \in Q_{1+AB}$ using SDP.

★ MDI min-entropic source

- we find zero min-entropy against Q_{1+AB}
- so we put further conditions on the observed statistics:
 $P(0, 0|l, l) = P(0, 0|m, m)$; $P(0, 0|l, l) = P(0, 0|m, m)$; $\forall l, m \in \{1, 2, 3, 4\}$
- interestingly positive min entropy is obtained for $I(P) < 0$



★ MDI min-entropic source

- two qubits entangled Werner class of states also satisfied the required conditions and hence are useful for MDI min-entropy (randomness) certification
- no separable state satisfies this condition $I(P) < 0$ hence no cheating strategy is possible by sharing separable correlations
- correlations in semi-quantum scenario cannot even be simulated by local operation and classical correlation (LOCC) (Rosset et al.)

★ Summary of the talk:

- Randomness is a useful resource
- Randomness certification is not possible **mathematically** and also in **classical world**
- In quantum world randomness generation is possible
- Bell's theorem: DI certification possible:– **nonlocal entangled states** are useful
- **local entangled states** are useful for MDI randomness certification

Thank You!